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Declarative Encodings of Acyclicity Properties

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Overview

Background

Encoding Directed Acyclic Graphs

Encoding Directed Forests

Encoding Directed Trees

Conclusions

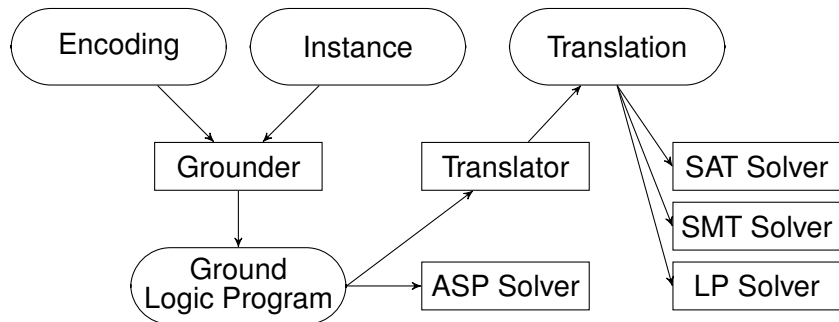
Motivation

- ▶ Numerous application problems involve the construction of **acyclic or tree structures**
 - ▶ Bayesian networks, Markov networks, Phylogenetic trees, ...
- ▶ Such structures are no basic primitives in common **constraint-based representation formalisms**, e.g.:
 - ▶ Answer Set Programming (ASP)
 - ▶ Boolean Satisfiability (SAT)
 - ▶ SAT Modulo Theories (SMT)
 - ▶ Mixed Integer Linear Programming (LP)

 Need for **compact and efficient encodings**

Approach

1. Uniform encoding in first-order ASP language
2. Automatic grounding w.r.t. instance data
3. Off-the-shelf solving in ASP or via translation to SAT / SMT / LP



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Directed Graph Representation

- ▶ The following **instance** represents graphs with
 - ▶ 5 nodes
 - ▶ complete set of edge candidates

```
node(1). node(2). node(3). node(4). node(5).  
pair(1,2). pair(1,3). pair(1,4). pair(1,5).  
pair(2,1). pair(2,3). pair(2,4). pair(2,5).  
pair(3,1). pair(3,2). pair(3,4). pair(3,5).  
pair(4,1). pair(4,2). pair(4,3). pair(4,5).  
pair(5,1). pair(5,2). pair(5,3). pair(5,4).
```

- ▶ Edge **generator** rule:

```
{ edge(X,Y) } :- pair(X,Y).
```

Acyclicity Testing

Acyclicity test procedure

While there is a leaf $v \in V$:

$$E := E \setminus \{\langle u, v \rangle \mid u \in V\}$$

$$V := V \setminus \{v\}$$

Return ($V \stackrel{?}{=} \emptyset$)

Declarative ASP encoding

`order(X,Y) :- pair(X,Y), not edge(X,Y).`

`order(X,Y) :- pair(X,Y), order(Y).`

`order(X) :- node(X), order(X,Y) : pair(X,Y).`

`:- node(X), not order(X).`

 Encoding is leaf-driven, non-tight, and linear

From Non-Tight to Tight (Leaf-Driven) Encoding

- ▶ Consider nodes with labels $1, \dots, n$

`order(X,Y,1..n) :- pair(X,Y), not edge(X,Y).`

`order(X,Y,N-1) :- pair(X,Y), order(Y,N), 1 < N.`

`order(X,N) :- node(X), order(X,Y,N) : pair(X,Y).`

`:- node(X), not order(X,1).`

- 👉 Encoding is leaf-driven, **tight**, and of size $\mathcal{O}(|E| \times |V|)$

Experimental Evaluation

600 edge candidates	leaf-driven encoding		root-driven encoding	
	non-tight	tight	non-tight	tight
clasp	0.31	80.26	0.34	0.39
clasp/sat	1.03	0.86	2.38	1.37
lingeling	2.62	1.78	1.82	1.70
z3	1.12	—	1.48	—
cplex	681.47	261.90	655.29	333.47

ASP solvers: clasp (3.1.0)

SAT solvers: clasp/sat, lingeling (ayv-86bf266-140429)

SMT solvers: z3 (4.3.2)

LP solvers: cplex (12.6.0.0)

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From Directed Acyclic Graphs to Forests

- ▶ Directed **forests** are directed acyclic graphs such that every node has **at most one** incoming edge
- ☞ Directed acyclic graph encodings can be augmented with respective **test**

Pairwise mutual exclusion

:- **edge**(X,Z) , **edge**(Y,Z) , X < Y.

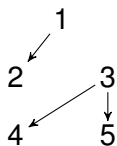
☞ Space complexity $\mathcal{O}(|E| \times |V|)$

Cardinality constraint

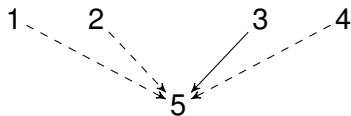
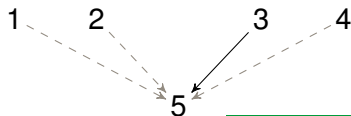
:- **node**(Y) , 2 { **edge**(X,Y) }.

☞ Space complexity $\mathcal{O}(|E|)$

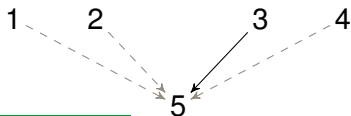
Linear “Normalized” At-Most-One Tests



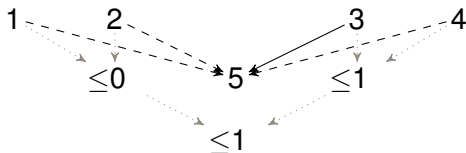
Linear traversal



Bidirectional traversal



Tournament traversal



Experimental Evaluation

clasp (3.1.0)	leaf-driven encoding		root-driven encoding	
	non-tight	tight	non-tight	tight
pairwise	5.93	8.18	7.44	8.95
cardinality	3.69	6.41	5.30	5.11
linear	7.29	8.47	8.50	7.08
bidirectional	4.68	6.44	6.14	6.94
tournament	12.36	8.56	8.41	8.61

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From Directed Forests to Trees

- ▶ Directed **trees** are directed forests such that
 - ▶ there is (exactly) **one root**
 - ▶ there are (exactly) $|V|-1$ **edges**
 - ▶ all nodes are **connected**

👉 Directed forest encodings can be augmented with **test(s)**

One root

`child(Y) :- edge(X,Y).`

+ any **at-most-one** encoding over instances of 'not child(Y)'

$|V|-1$ edges

`:- not n-1 { edge(X,Y) } n-1.`

Connectedness

`reach(X) :- reach(Y), edge(X,Y). reach(1).`

`reach(Y) :- reach(X), edge(X,Y). :- node(X), not reach(X).`

Experimental Evaluation

clasp (3.1.0)	leaf-driven encoding		root-driven encoding	
	non-tight	tight	non-tight	tight
pairwise	2.29	3.73	2.78	3.02
cardinality	2.14	4.09	2.77	5.38
linear	8.42	11.06	8.93	46.92
bidirectional	5.46	6.47	2.87	3.99
tournament	11.56	5.09	11.68	7.81
$\geq V -1$ edges	10.16	28.25	15.86	7.82
$= V -1$ edges	12.81	16.43	12.18	38.47
connectedness	2.30	5.85	2.72	46.07

Conclusions

- ▶ Acyclic or tree structures are central in **many applications**
- ▶ Need for development and study of **declarative encodings**
- ▶ First-order ASP language (with recursion, cardinality and weight constraints, etc.) facilitates **exploration** of encodings
- ▶ Diverse formulations of acyclicity, forest, and tree conditions yield rich family of encoding **variants** to experiment with
- ▶ Grounding and automatic translations provide large variety of **back-end solvers**
- ▶ Given encodings furnish **templates** for corresponding tasks
- ▶ See paper for further study of logic-based characterizations of **undirected** forests and trees as well as chordal graphs