Computer Formalization \ of the Real Numbers in Lean 4 **EROFOTALL** HIIT Foundations Friday The Aalto Formalization Alliance Niklas Halonen 2025-09-26 Aalto University



$$\mathbb{N}\stackrel{?}{\subset}\mathbb{Z}\stackrel{?}{\subset}\mathbb{Q}\stackrel{?}{\subset}\mathbb{R}$$

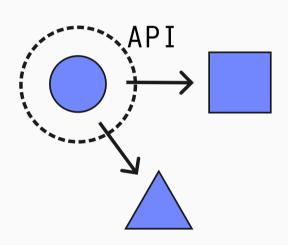


$$\mathbb{N}\stackrel{?}{\subset}\mathbb{Z}\stackrel{?}{\subset}\mathbb{Q}\stackrel{?}{\subset}\mathbb{R}$$

Semigroup \supseteq Monoid \supseteq Group \supseteq Ring \supseteq Field

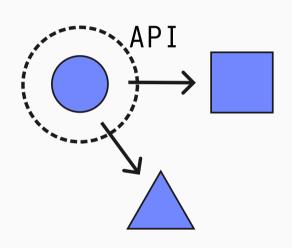


- API (Application Programming Interface) is the interface that is exposed by an object.
 - In **mathematics**: the object's *relationship* to other objects.
 - In **computer science**: the object's constructors, field accessors, methods.



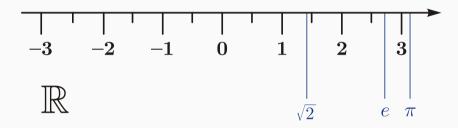


- API (Application Programming Interface) is the interface that is exposed by an object.
 - In **mathematics**: the object's *relationship* to other objects.
 - In **computer science**: the object's constructors, field accessors, methods.
- Shared behavior can be organized in type classes.
- In Mathlib (Lean's mathematics library) mathematical properties are expressed using a **type class hierarchy**.
 - Examples: CommRing, Field, MetricSpace.





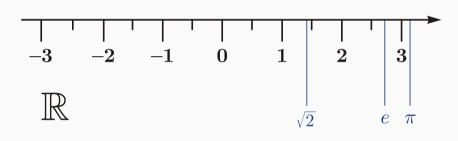
Informally



- The number line
- "Rationals ∪ irrationals"
- Convergent sequences have a limit



Informally

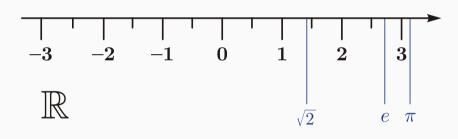


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- R is a "conditionally complete linear ordered field"
 - Conditionally complete: every nonempty subset of R bounded from above has a least upper bound (in R).



Informally

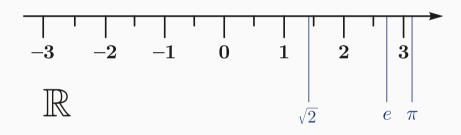


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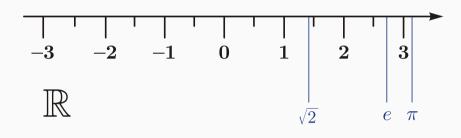


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- Any such R is isomorphic with \mathbb{R} .

What is Computer Formalization?











Organization

Verification

Automation



Organization

Create **libraries** of intercompatible definitions and theorems.

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Proofs are systematically verified by a **trusted kernel**.

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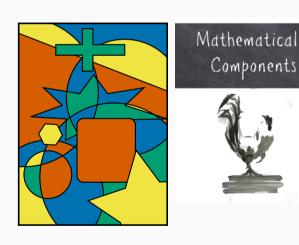
LLMs can verifiably prove theorems that **humans** care about.

History of Computer Formalization



2005 - 2012

Four-color and Feit-Thompson theorem

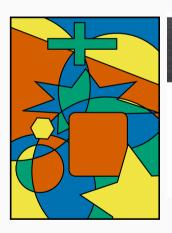


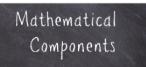
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2014

Kepler conjecture

The Flyspeck project has been a large team effort over a period of years, and my contributions have been just a fraction of the whole. This appendix cites their contributions. The principal formalizer of each chapter is indicated with an asterisk.

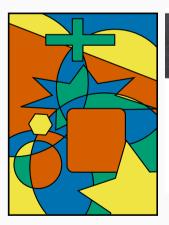
Text Chapter	Author of formalization work
Trigonometry	Nguyen Quang Truong*; Rute, Jason;
	Harrison, John; Vu Khac Ky
Volume	Harrison, John*; Nguyen Tat Thang
Hypermap	Tran Nam Trung*
Fan	Hoang Le Truong*; Harrison, John
Packing	Solovyev, Alexey*; Vu Khac Ky*;
	Nguyen Tat Thang; Hales, Thomas
Local Fan	Nguyen Quang Truong*; Hoang Le Truong*
Tame Hypermap	Solovyev, Alexey*; Dat Tat Dang; Trieu Thi Diep;
	Vu Quang Thanh; Vuong Anh Quyen
Code-Verification	Author of formalization work
Hypermap Generation	Nipkow, Tobias*; Bauer, Gertrud*
Hypermap Generation Linear Programs	Nipkow, Tobias*; Bauer, Gertrud* Obua, Stephen*; Solovyev, Alexey*
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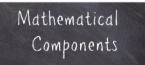
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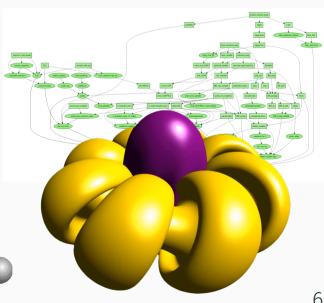
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Hypermap Generation Nipkow, Tobias*; Bauer, Gertrud* Linear Programs Obua, Stephen*; Solovyev, Alexey

Nonlinear Inequalities Solovyev, Alexey*

2019 -

Perfectoid Spaces, Liquid Tensor Experiment, Sphere Eversion, Mathlib



Formal Definition of the Reals





JSE says:

December 11, 2020 at 10:24 pm

Thanks for a very informative and idea-rich post, Peter!

Kevin, how are real numbers formalized in Lean in the first place? Via Cauchy sequences? If so, the idea of proving something about real vector spaces by exhibiting R as a quotient of a subring $Z[[T]]_{>r}$ of Z[[T]] cut out by archimedean norm conditions does somehow seem in the spirit of things; indeed I suppose you could think of this description of R as an alternate means of formalizing the real numbers, one particularly well-suited for functional analysis!



Reply



xenaproject says:

December 11, 2020 at 10:55 pm

of R as an alternate means of formalizing the real numbers, one particularly well-suited for functional analysis!



★ Like

Reply



 $\underline{\mathbf{xenaproject}}\ says: \leftarrow \mathsf{Kevin}\ \mathsf{Buzzard}$

December 11, 2020 at 10:55 pm

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★ Like

Reply



<u>xenaproject</u> says: ← Kevin Buzzard

December 11, 2020 at 10:55 pm

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★ Like

<u>Reply</u>



<u>xenaproject</u> says:

Kevin Buzzard

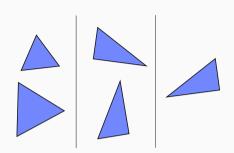
December 11, 2020 at 10:55 pm

Quotients



Examples of quotients from mathematics:

- Congruent triangles
- Modular arithmetic, $\mathbb{Z}/3\mathbb{Z}$
- Real numbers, $\{f: \mathbb{N} \to \mathbb{Q} \mid f \text{ Cauchy}\}/\approx$
 - Where $f \approx g$ if they are "eventually ε -close"



Equivalence classes of congruent triangles

Quotients

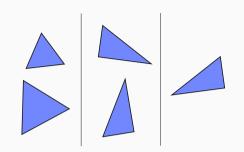


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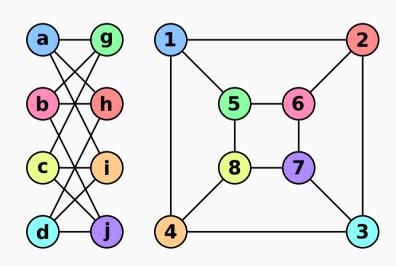
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Examples of quotients from computer science:

- Multisets, a.k.a. lists modulo permutation
- Isomorphism classes of graphs: e.g. labeled graphs modulo relabeling of vertices



Equivalence classes of congruent triangles

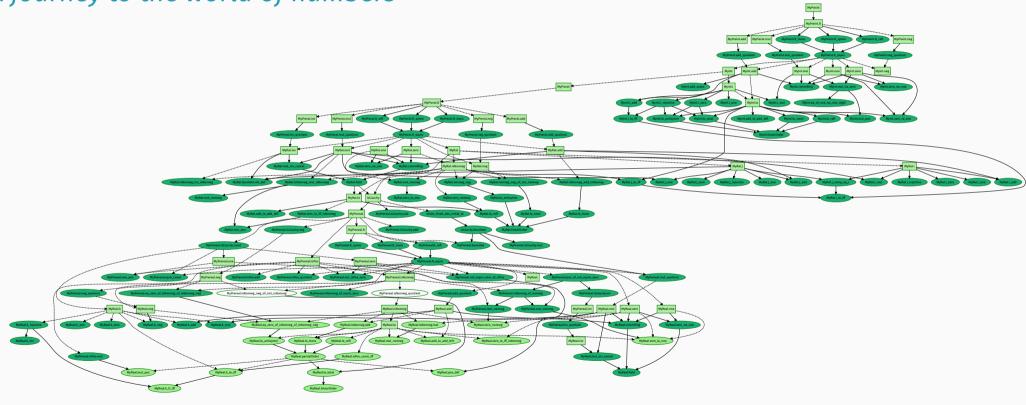


Isomorphic graphs

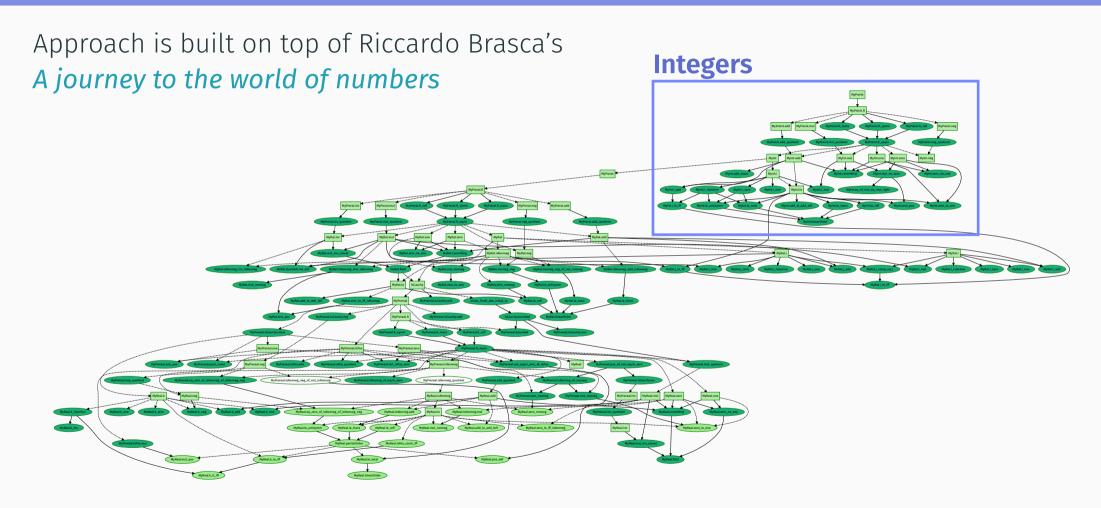


Approach is built on top of Riccardo Brasca's

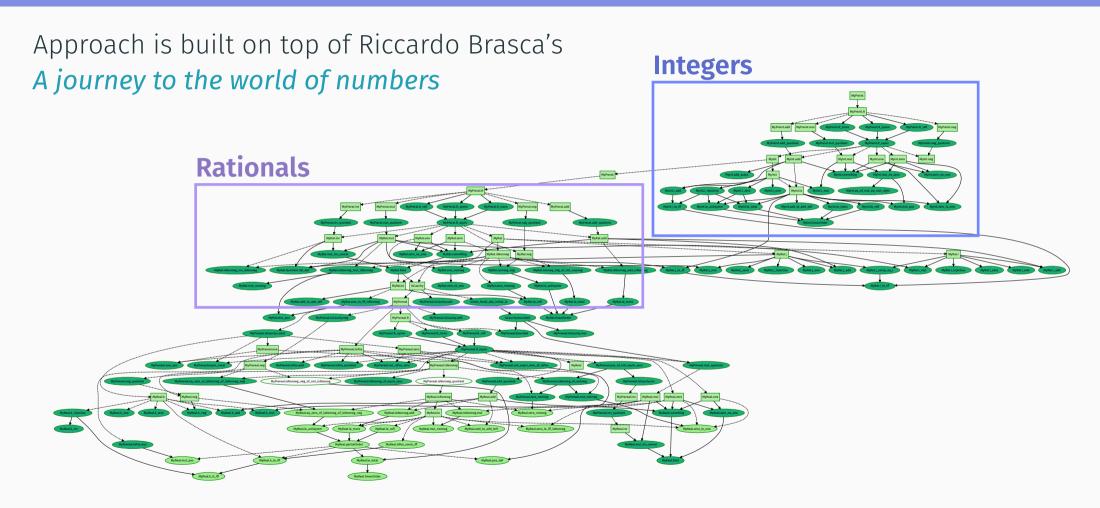
A journey to the world of numbers



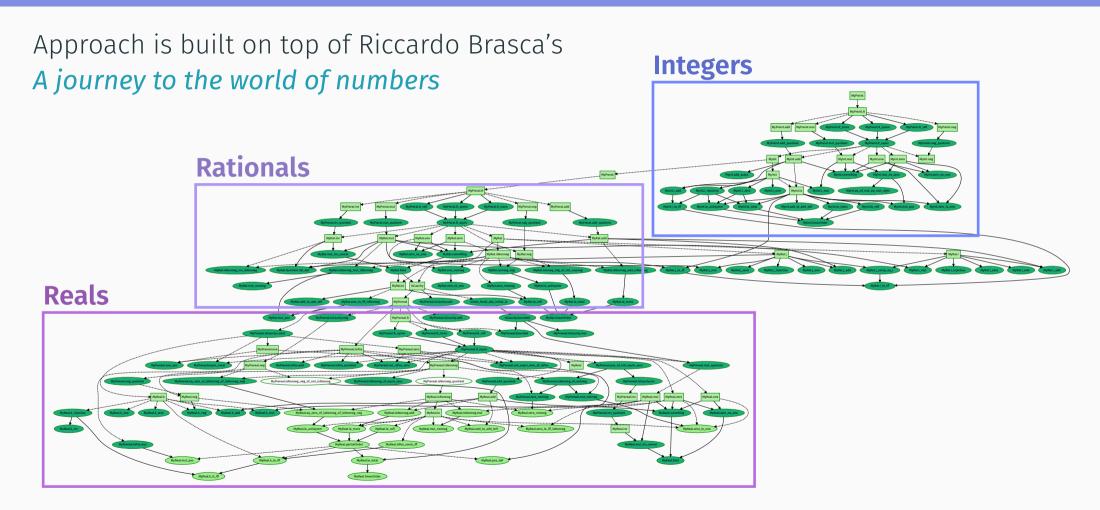


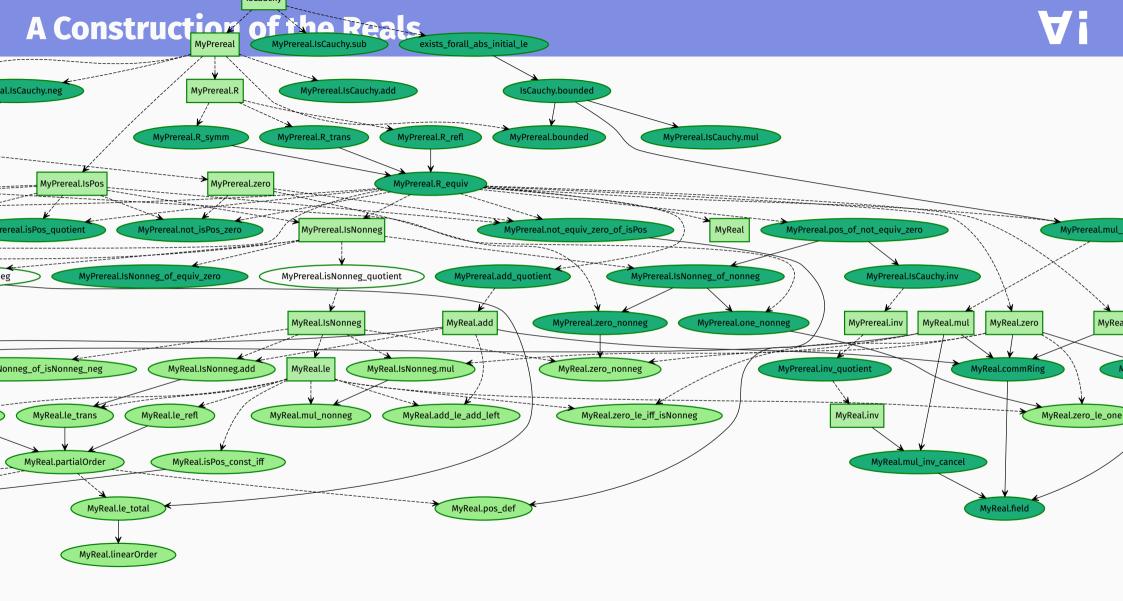


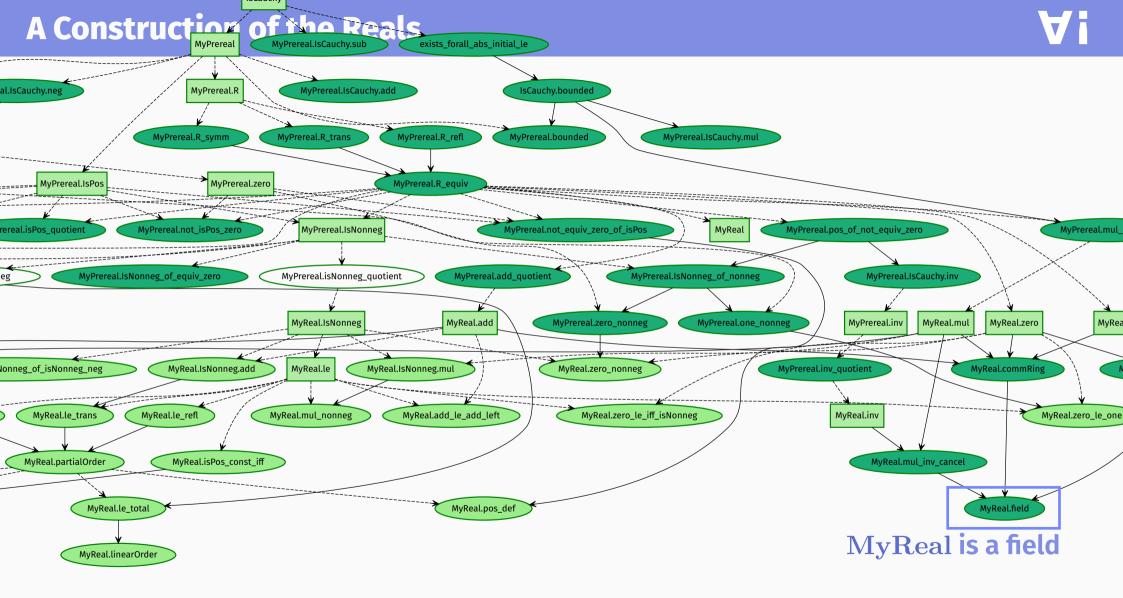




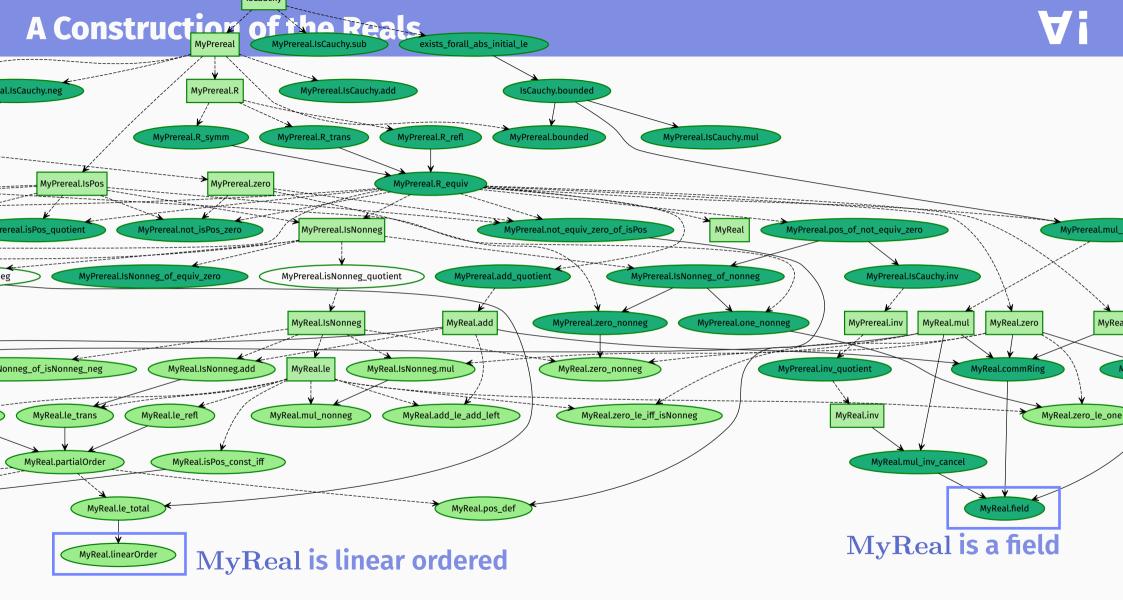








Niklas Halonen, Aalto University



Niklas Halonen, Aalto University

11 / 19

$MyReal \simeq \mathbb{R}$



- We can prove that our real numbers are the real numbers.
 - Mathlib gives the proof of isomorphism between all conditionally complete linear ordered fields:

LinearOrderedField.inducedOrderRingIso MyReal \mathbb{R} : MyReal $\simeq +*o$ \mathbb{R}

- Some API for real numbers, such as metric spaces (distances) is not immediately available for MyReal.
 - It is technically possible to use this for constructing the metric space structure, but I found it difficult to use in proofs of continuity for instance.

Metaprogramming and Tactics

Short Proof: Associativity of Addition



x = [a]

Let x, y, z: MyReal, i.e. equivalence classes of Cauchy sequences. Addition is associative: (x + y) + z = x + (y + z).

To show:

$$\vdash (\llbracket a \rrbracket + \llbracket b \rrbracket) + \llbracket c \rrbracket = \llbracket a \rrbracket + (\llbracket b \rrbracket + \llbracket c \rrbracket).$$
 a is a Cauchy sequence

Proof: The equivalence classes are equal if their representatives are related.¹

$$\vdash (a+b)+c\approx a+(b+c)$$

$$\vdash \forall \varepsilon>0, \exists N, \forall n\geq N, |(a_n+b_n)+c_n-(a_n+(b_n+c_n))|\leq \varepsilon$$

Let $\varepsilon > 0$, use N := 0 and let $n \ge N$. Simplifying the left-hand side using ring axioms on rationals reduces the goal to $|0| \le \varepsilon$, which follows from $\varepsilon > 0$.

¹This is called the *quotient axiom* in Lean.

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The other commutative ring axioms hold by a "similar argument".

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Metaprogramming



```
macro "quot_proof" : tactic =>
 `(tactic|
  focus
    refine Quot.ind fun a => ?_
   try refine Quot.ind fun b => ?_
    try refine Quot.ind fun c => ?
    apply Quot.sound
    intro ε hε
    simp
    use 0
    intro n hn
   try {ring_nf; simp [hɛ.le]}
```

Metaprogramming



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    intro \varepsilon he \leftarrow let \varepsilon > 0
    \mathsf{use} \ \mathbf{0} \qquad \qquad \longleftarrow \mathsf{use} \ N \coloneqq 0
     intro n hn \longleftarrow let n \geq N
     try {ring_nf; simp [hε.le]}
              simplify with ring axioms
```

Metaprogramming



```
instance commRing : CommRing MyReal where
macro "quot proof" : tactic =>
                                                  add := (\cdot + \cdot)
  `(tactic|
                                                  add assoc := by quot proof
  focus
                                                  zero := 0
    refine Quot.ind fun a => ?
                                                  zero add := by quot proof
    try refine Quot.ind fun b => ?_
                                                  add zero := by quot proof
    try refine Quot.ind fun c => ?
                                                  add comm := by quot proof
    apply Quot.sound
                                                  mul := (\cdot * \cdot)
    intro \varepsilon he \leftarrow let \varepsilon > 0
                                                  left distrib := by quot proof
    simp ← simplify
                                                  right distrib := by quot proof
    \mathsf{use} \ \mathbf{0} \qquad \qquad \longleftarrow \mathsf{use} \ N := 0
                                                  zero_mul := by quot_proof
                                                  mul zero := by quot proof
    intro n hn \longleftarrow let n > N
                                                  mul assoc := by quot proof
    try {ring nf; simp [hɛ.le]}
                                                  one := 1
             simplify with ring axioms
                                                  one mul := by quot proof
                                                  mul one := by quot proof
                                                  neg := (-\cdot)
                                                  mul comm := by quot proof
```

neg add cancel := by quot proof

Recap & Key points



- 1. API and type classes
 - It's useful to group shared behavior in type classes, even when formalizing mathematics. The theory can be generic (agnostic) over the construction.
 - However, concrete constructions are often easier to work with.

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 - We can what real numbers are in code, and get all of their properties.
 - Key "characteristic" properties can be expressed using type classes.
 - The existence of a construction proves that the theory of real numbers is consistent.

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- The existence of a construction proves that the theory of real numbers is consistent.

3. Metaprogramming

- Proof that real numbers are a field boils down to field axioms on rationals.
 - We can make "cases are similar" arguments exact with computer formalization by defining a custom tactic.

Thank you for listening!

Questions?

Join the Aalto Formalization Alliance (ForAlli) for discussion.

Want to learn Lean? See these learning resources.

To contact me, send email to niklas. 2. halonen AT aalto. fi.

Learning Resources From the Community Page



Learning Resource	Туре	Audience	Level	Requirements
Natural Number Game	Exercises (game)	General	High-school	In-browser only
Lean game server (including NNG)	Exercises (game)	General	High-school	In-browser only
Mechanics of Proof	Book + Exercises	General	Undergraduate	Browser or local
GlimpseOfLean	Exercises	Math	Undergraduate math (1st year)	Browser is sufficient
Introduction to the Lean 4 theorem prover and programming language by Leonardo de Moura	Video	CS	Undergraduate CS (1st year)	YouTube
Lean Documentation Overview	Book	CS	Undergraduate CS	
Functional Programming in Lean (FPIL)	Book + Exercises	CS	Undergraduate CS (2nd year)	Local installation
Theorem Proving in Lean (TPIL)	Book + Exercises	CS	Undergraduate CS (3rd year)	Local installation
CS 99: Functional Programming and Theorem Proving in Lean 4	Course Slides + Exercises	CS + Math	Undergraduate and graduate	Local installation
Mathematics in Lean (MIL) mathematics_in_lean	Book + Exercises	Math	Undergraduate and graduate math	Local installation
Formalising Mathematics Handbook formalising-mathematics-notes	Notes + Exercises	Math	Undergraduate and graduate math	Local installation
The Hitchhiker's Guide to Logical Verification	Book + Exercises	CS	Undergraduate and graduate CS	Local installation

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